



Friedrich-Alexander-Universität Erlangen-Nürnberg

Introduction to Parallel Programming with MPI

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Blocking Collective Communication



Collectives: operations including all ranks of a communicator

All ranks must call the function!

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- Blocking variants: buffer can be reused after return
- Nonblocking variants (since MPI 3.0): buffer can be used after completion (MPI_Wait*/MPI_Test*)
- May or may not synchronize the processes
- Cannot interfere with point-to-point communication
 - Completely separate modes of operation!

- Rules for all collectives
 - Data type matching
 - No tags
 - Count must be exact, i.e., there is only one message length, buffer must be large enough

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 - Collective computation (reduction, scan)
 - Combinations of data movement and computation (reduction + broadcast)

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 - Collective computation (reduction, scan)
 - Combinations of data movement and computation (reduction + broadcast)
- General assumption: MPI does a better job at collectives than you trying to emulate them with a collection of point-to-point calls

Barrier

 Explicit synchronization of all ranks from specified communicator

```
MPI_Barrier(comm);
```

- Ranks only return from call after every rank has called the function
- MPI_Barrier: rarely needed
 - Debugging

Barrier

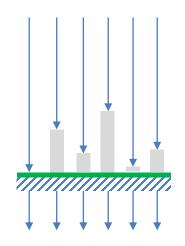
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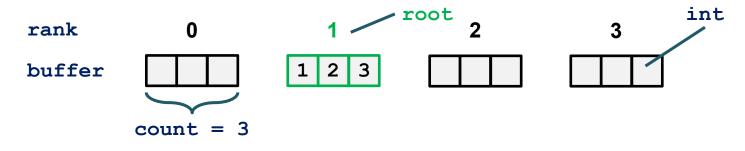
```
  Debugging
```



Send buffer contents from one rank ("root") to all ranks

MPI_Bcast(buf, count, datatype, int root, comm);

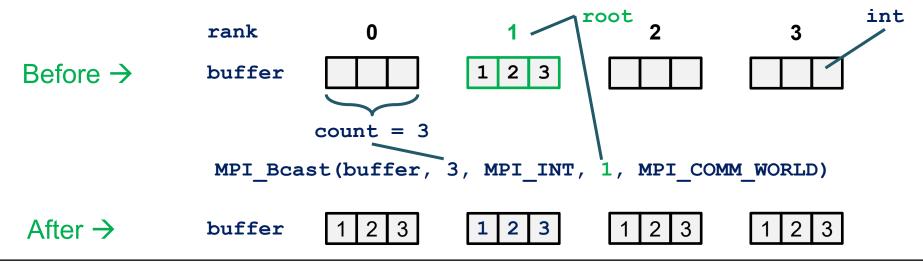
no restrictions on which rank is root – often rank 0

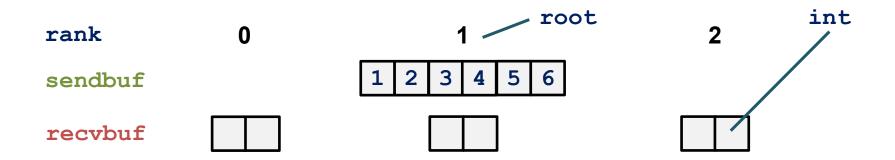


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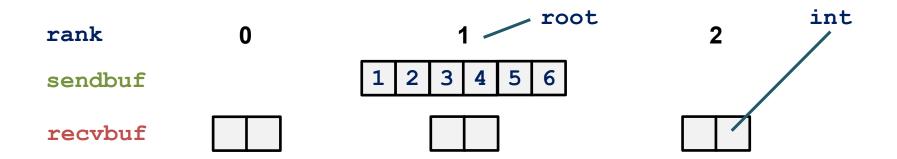
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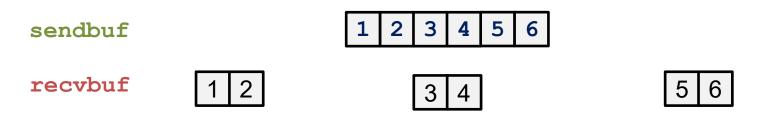




MPI_Scatter(sendbuf, 2, MPI_INT, recvbuf, 2, MPI_INT, root, MPI_COMM_WORLD)



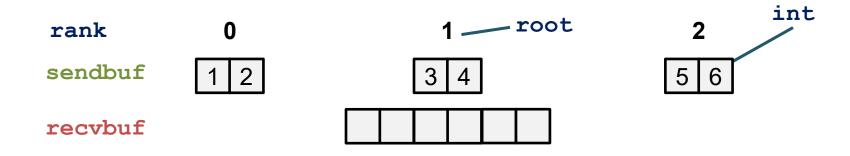
MPI_Scatter(sendbuf, 2, MPI_INT, recvbuf, 2, MPI_INT, root, MPI_COMM_WORLD)



Send every i-th chunk of an array to the i-th rank

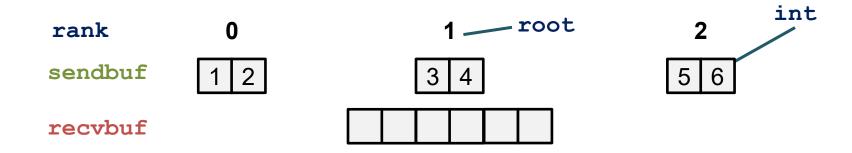
- Root and comm must be the same on all processes
- Type signature of send and receive variables must match
- Usually, sendcount = recvcount because sendtype = recvtype
 - This is the length of the chunk
- sendbuf is ignored on non-root ranks because there is nothing to send

Gather

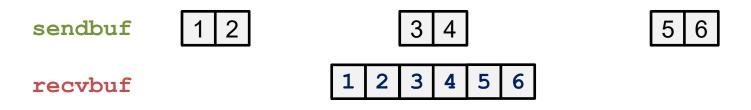


MPI_Gather(sendbuf, 2, MPI_INT, recvbuf, 2, MPI_INT, root, MPI_COMM_WORLD)

Gather



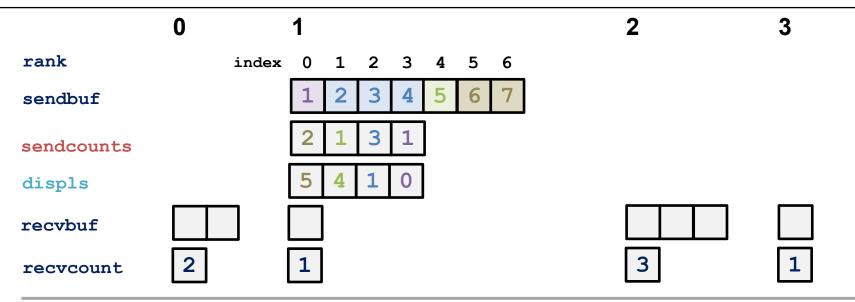
MPI_Gather(sendbuf, 2, MPI_INT, recvbuf, 2, MPI_INT, root, MPI_COMM_WORLD)



 Receive a message from each rank and place i-th rank's message at i-th position in receive buffer

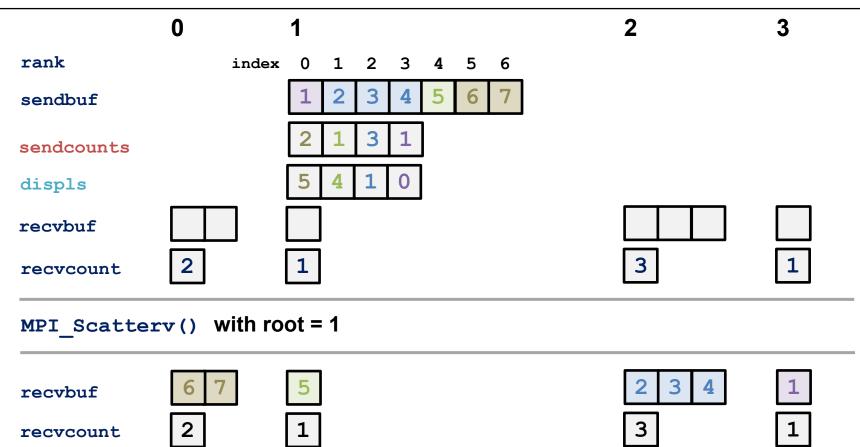
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Scatterv



MPI_Scatterv() with root = 1

Scatterv



Scatterv: more flexible scatter

Send chunks of different sizes to different ranks

```
MPI_Scatterv(
  sendbuf, int sendcounts[], int displs[], sendtype,
  recvbuf, recvcount, recvtype, root, comm)
```

sendcounts[]: array specifying the number of elements to send to
 each rank: send sendcounts[i] elements to rank i

displs[]: integer array specifying the displacements in
 sendbuf from which to take the outgoing data to
 each rank, specified in number of elements

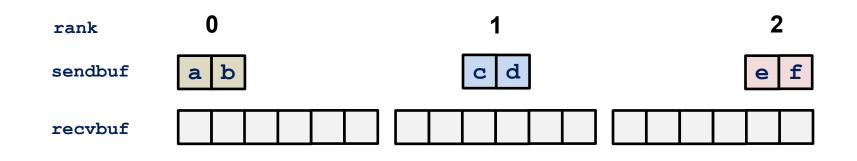
Gatherv: more flexible gather

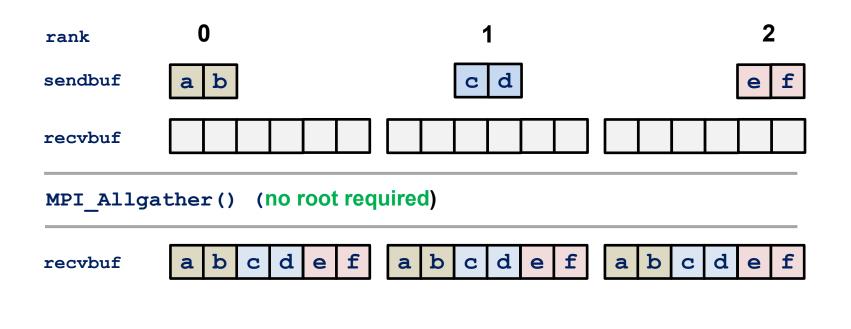
Receive segments of different sizes from different ranks

```
MPI_Gatherv(
  sendbuf, sendcount, sendtype,
  recvbuf, int recvcounts[], int displs[], recvtype,
  root, comm)
```

recvcounts[]: array specifying the number of elements to receive
from each rank: receive recvcounts[i] elements from rank i

displs[]: integer array specifying the displacements where received data from specific rank is put in recvbuf, in units of elements:



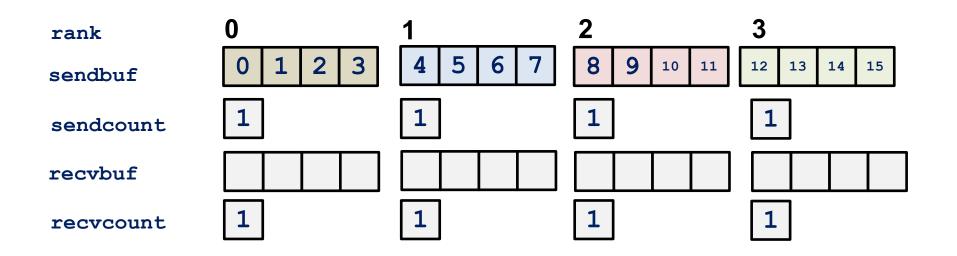


In this example: sendcount=recvcount=2

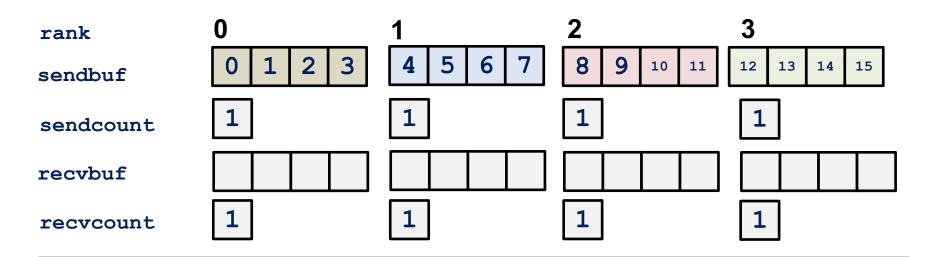
Combination of gather and broadcast

- Also available: MPI_Allgatherv() (cf. MPI_Gatherv())
- Why not just use gather followed by a broadcast instead?
 - MPI library has more options for optimization
 - General assumption: Combined collectives are faster than using separate ones

Alltoall

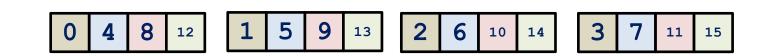


Alltoall



MPI_Alltoall() (no root required)

recvbuf



MPI_Alltoall: For all ranks, send i-th chunk to i-th rank

- MPI_Alltoallv: Allows different number of elements to be send/received by each rank
- MPI_Alltoallw: Allows also different data types and displacements in bytes

Summary of MPI Collective Communications

- MPI (blocking) collectives
 - All ranks in communicator must call the function
- Communication and synchronization
 - Barrier, broadcast, scatter, gather, and combinations thereof
- In-place buffer specification MPI_IN_PLACE
 - Save some space if needed

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Answer: b., it is necessary.

- 4. To send an identical piece of data to all other processes in a communicator, which collective call should be used?
 - a. MPI_Gather
 - b. MPI_Bcast
 - c. MPI_Scatter
 - d. MPI_Alltoall

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Answer: d.