



Introduction to Parallel Programming with MPI

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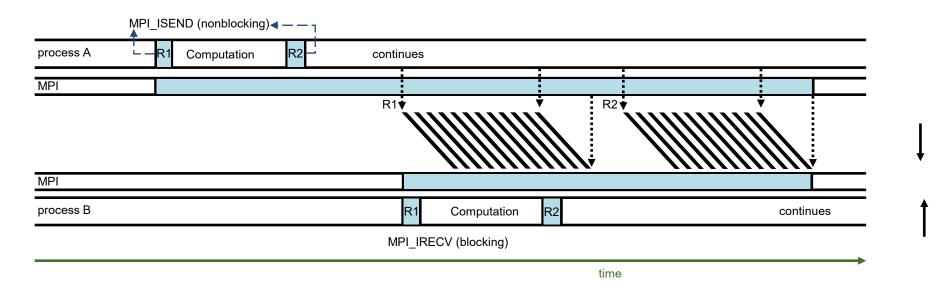
MPI Point-to-Point Communication: Nonblocking



Nonblocking point-to-point communication

- Call to a nonblocking send/recv procedure returns straight away. It avoids synchronization so that the following opportunities can be exploited:
 - Avoiding certain deadlocks
 - Truly bidirectional commun.

- Avoid idle time:
 - Overlapping commun. and comput.



Standard nonblocking send/receive

```
MPI_Isend(sendbuf, count, datatype, dest, comm, MPI_Request * request);

MPI_Irecv(recvbuf, count, datatype, source, tag, comm, MPI_Request * request);

request:pointer to variable of type MPI_Request, will be associated with the corresponding operation
```

- Do not reuse sendbuf/recvbuf before MPI_Isend/MPI_Irecv has been completed
 - Return of call does not imply completion
- MPI Irecv has no status argument
 - obtained later during completion via MPI Wait*/MPI Test*

Nonblocking send and receive variants

Completion

- Return of MPI I* call does not imply completion
- Check for completion via MPI_Wait* / MPI_Test*
- Semantics identical to blocking call after successful completion

nonblocking MPI function	blocking MPI function	type	completes when
MPI_Isend	MPI_Send	synchronous or buffered	depends on type
MPI_Ibsend	MPI_Bsend	buffered	buffer has been copied
MPI_Issend	MPI_Ssend	synchronous	remote starts receive
MPI_Irecv	MPI_Recv		message was received

Test for completion

Two test modes:

- Blocking
 - MPI_Wait*: Wait until the communication has been completed and buffer can safely be reused
- Nonblocking
 - MPI_Test*: Return true (false) if the communication has (not) completed

Despite the naming, the modes both pertain to nonblocking point-to-point communication!

Test for completion – single request

Test one communication handle for completion:

```
MPI Wait(MPI Request * request,
           MPI Status * status);
MPI Test(MPI Request * request, int * flag,
           MPI Status * status);
request: request handle of type MPI Request
status: status object of type MPI Status (cf. MPI Recv)
flag: variable of type int to test for success
```

Use of wait/test

MPI_Wait

```
MPI Request request;
MPI Status status;
MPI Isend(
  send buffer, count, MPI CHAR,
  dst, 0, MPI COMM WORLD, &request);
// do some work...
// do not use send buffer
MPI Wait(&request, &status);
// use send buffer
```

MPI_Test

```
MPI Request request;
MPI Status status;
int flag;
MPI Isend(
  send buffer, count, MPI CHAR,
  dst, 0, MPI COMM WORLD, &request);
do {
    // do some work...
    // do not use send buffer
    MPI Test(&request, &flag, &status);
} while (!flag);
// use send buffer
```

Wait for completion – all requests in a list

- MPI can handle multiple communication requests
- Wait/Test for completion of multiple requests:

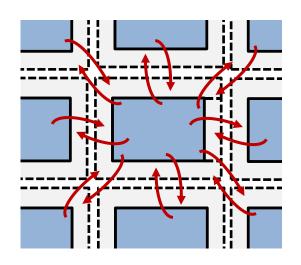
Waits for/Tests if all provided requests have been completed

Use of MPI Waitall

```
Arrays of
MPI Request requests[2];
                                               requests and
MPI Status statuses[2];
                                                 statuses
MPI Isend(send buffer, ..., &(requests[0]));
MPI Irecv(recv buffer, ..., &(requests[1]));
                                 number of elements in
// do some work...
                                     the arrays
MPI Waitall(2, requests, statuses)
// Isend & Irecv have been completed
```

Ghost cell exchange with nonblocking MPI

Ghost cell exchange with nonblocking send/recv with all neighbors at once



Possible implementation:

- 1. Update cells that need the halo
- 2. Copy new data into contiguous send buffers
- 3. Start nonblocking receives/sends from/to corresponding neighbors
- 4. Update local cells that do not need halo cells for boundary conditions ("bulk update")
- 5. Wait with MPI_Waitall for all obtained requests to complete
- 6. Copy received halo data into ghost cells

Opportunity to overlap communication (steps 3-5) with bulk update (MPI implementation permitting)

Wait for completion – one or several requests out of a list

Wait for/Test if exactly one request among many has been completed

Wait for/Test if at least one request among many has been completed

Use of MPI Testany

```
MPI Request requests[2];
MPI Status status;
int finished = 0;
                                                  completed requests are
MPI Isend(send buffer, ..., &(requests[0]));
                                                   automatically set to
MPI Irecv(recv buffer, ..., &(requests[1]));
                                                   MPI REQUEST NULL
                                                  completed requests:
do {
                                                   requests[idx]
  // do some work...
  MPI Testany(2, requests, &idx, &flag, &status);
  if (flag) { ++finished; }
} while (finished < 2);</pre>
```

Pitfalls with nonblocking MPI and compiler optimizations

Fortran:

```
MPI_IRECV(recvbuf, ..., request, ierror)
MPI_WAIT(request, status, ierror)
write (*,*) recvbuf
```

may be compiled as
MPI_IRECV(recvbuf, ..., request, ierror)
registerA = recvbuf
MPI_WAIT(request, status, ierror)
write (*,*) registerA

MPI might modify recybuf after MPI_IRECV returns, but the compiler has no idea about this

- i.e., old data is written instead of received data!
- Workarounds:
 - recvbuf may be allocated in a common block, or
 - calling MPI_GET_ADDRESS(recvbuf, iaddr_dummy, ierror)
 after MPI WAIT
 - asynchronous attribute

Nonblocking point-to-point communication

- Standard nonblocking send/recv MPI_Isend()/MPI_Irecv()
 - Return of call does not imply completion of operation
 - Use MPI_Wait*() / MPI_Test*() to check for completion using request handles
- All outstanding requests must be completed!
- Potentials
 - Overlapping of communication with work (not guaranteed by MPI standard)
 - Overlapping send and receive
 - Avoiding synchronization and reducing idle times
- Caveat: Compiler does not know about asynchronous modification of data

Quiz:

- Every nonblocking send or receive requires a subsequent MPI_Wait* or MPI_Test* call?
 - a) Correct

- b) Incorrect
- 2. Can MPI_Isend be matched with blocking receive (MPI_Recv)?
 - a) Yes

- b) No
- 3. Which one is not a potential use case of nonblocking point-to-point communication in MPI?
 - a) Overlapping send and receive
 - b) Avoiding every possible idle times
 - c) Overlapping of communication with work